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# Attack Class: Address Spoofing

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23 Oct 1996

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# Overview of Talk

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- Introduction
- Background material
- Attack class
- Example attack
- Popular questions
- Extensions

# UCD Vulnerabilities Group

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- UCD's vulnerabilities group studies attacks and their underlying vulnerabilities for the purpose of modeling them. We believe a sufficiently complete model will allow us to both predict new instances of general attack classes and build generic schemes for detecting exploitations of general vulnerability classes.

# Address Masquerading

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- Many of today's network services use host names or addresses for both identification AND authentication.
- Examples: rlogin, rsh, mountd, wrappers, firewalls
- Higher level services use these lower level services (e.g., backups)

# History of Talk

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- R.T. Morris, 85
- S. Bellovin, 89
- UCD Discussed, Feb. 94
- UCD Presented, Mar 94
- Mitnick-Tsutomu, Dec 94
- UCD paper, spring 95
- Mendax, Rbone, summer 95
- Wee (UCD), fall 95
- USAF project, Jan. 96

# Orders and Dialogues

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- Need better names
  - » asynchronous vs. synchronous
  - » connectionless vs. connection-oriented
- An order is a request requiring only a single “message”.
- A dialogue is a request which requires the exchange of several, interdependent “messages”.
- From recipient’s point of view

# Connectionless Communication (Orders)

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- Connectionless communication (e.g., supplied by UDP), does not keep state information
- No guarantee of delivery or order
- Efficient in many environments
- RPC on UDP (NFS)

# Connection-oriented Communication (Dialogues)

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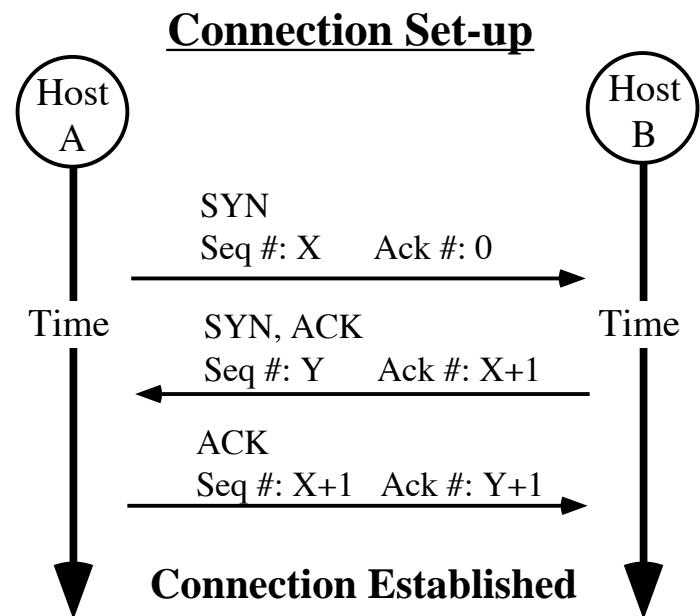
- Additional state information kept, representing a limited history of communication
- Provides “guarantee” that information will both arrive and arrive in order
- May require more resources and be less efficient in some environments



# TCP / IP Example

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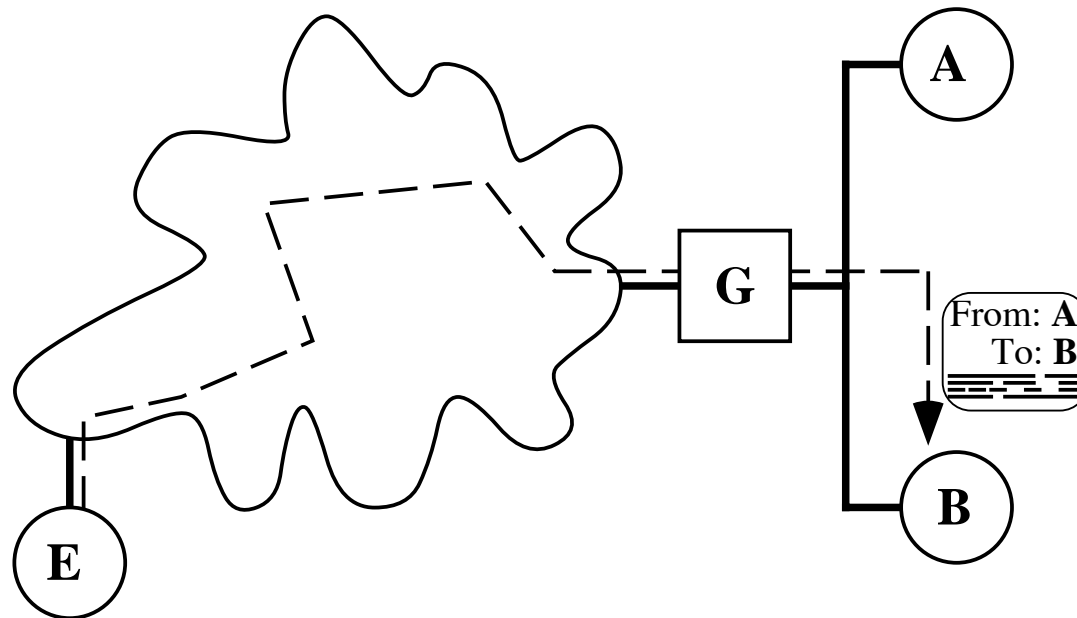
- Three phases: set-up, data exchange, tear-down
- set-up is a three-way handshake
- Third packet requires information from second packet.



# Routing in an internet

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- Host constructs packet and simply places it on the network
- As the packet travels across the internet, only the destination address is used

# The Attack

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- Definition of what an attack is
- Restrictions to be concerned with
- Strategy of the attacker

# Definition of Attack

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- Players: Alice (A), Bob (B), and Eve (E)
- Bob grants Alice special privileges by listing Alice's address or name in a special file
- Eve is the villain
- Eve's goal: **To get Bob to perform a specific action that he would perform for Alice but not Eve**

# Restrictions

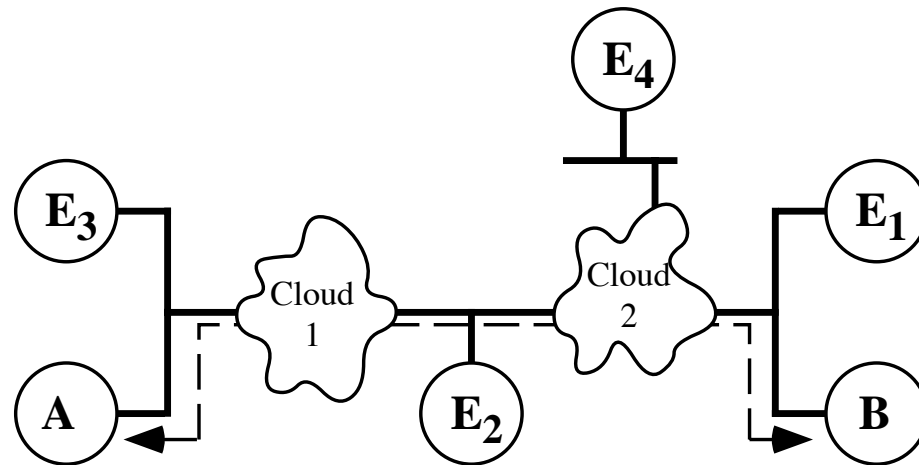
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- The placement of Alice, Bob, and Eve (the topology)
- The nature of the communication required by Eve to carry out the attack.
- These restrictions will help define Eve's strategy

# Architecture (or Topology)

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- Alice and Bob on separate networks; Eve in one of four locations
- Other architectures are simply special cases of this one

# Nature of Communication

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- Eve's communication must be indistinguishable from Alice's communication with Bob
- Order communication
  - » request is carried out immediately
  - » No role-backs
- Dialogue communication
  - » must make sense to Bob
  - » Alice cannot be allowed to interfere

# Eve's Strategy

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- Establish a forged communication with Bob
- Prevent Alice from alerting Bob until it is too late



# Establishing a Forged Communication

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- Construct packet, and place it on the network. The network will deliver it for Eve
- For order-based communication, the communication is done
- For dialogue-based communication, further messages must be exchanged
  - » if Eve is in  $E_1$ ,  $E_2$ , or  $E_3$ , further communication is easy
  - » if Eve is in  $E_4$ , she must either modify the messages' routes, or predict what the messages will contain

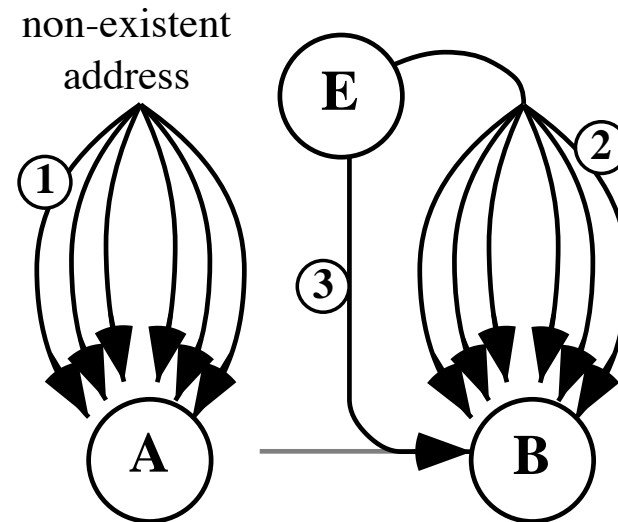
# Prevent Alice from Interfering

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- Prevent Bob's packets from reaching Alice (or Alice's from reaching Bob)
- Take away Alice's ability to respond
  - » wait for Alice to go down for maintenance
  - » force Alice to crash
  - » block part of Alice's operating system from processing Bob's packets (graceful ??)
- Complete communication before Alice can respond

# Example Attack

<b>Players</b>	<b>E</b>	adversary
	<b>A</b>	server
	<b>B</b>	X-client
<b>Steps</b>	<b>1</b>	Prevent Alice From Responding
	<b>2</b>	Probe for sequence number prediction
	<b>3</b>	Forge communication

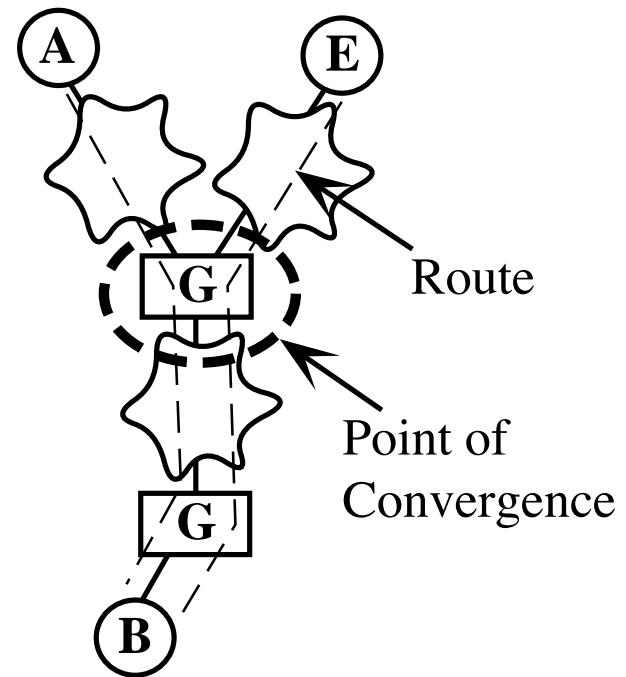


- Used against Tsutomu Shimamura, attributed to Kevin Mitnick
- Detailed ten years earlier by R.T. Morris

# Questions

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- Couldn't this attack be stopped by simply configuring routers not to forward obviously forged packets?





# Questions cont.

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- **Couldn't we simply write a more secure algorithm for choosing initial sequence numbers?**
- Only if Eve is NOT in position  $E_1$ ,  $E_2$ , or  $E_3$ , and Eve is NOT able to alter the path of Bob's messages to Alice (e.g., source routing or routing table modification). Also, this solution does not apply to order-based communications.

# Extensions to this Attack: Session Hijacking

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- One-time authentication services are vulnerable
- Commercial programs exist which do session hijacking
- Demonstrated against systems with challenge-response authentication

# Extensions cont.

- Eve's goal: **To get Bob to accept information he would only accept from Alice**

