# Modeling Vulnerabilities from buffer overflows to insider threat

#### **Sophie Engle NSF I/UCRC CIP MEETING**

UC Davis Kemper Hall 1008 · Tuesday June 17 2008

### Motivation

What does it mean for a system to be secure?

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### **Motivation**

What does it mean for a system to be secure?

physically secure?

What does it mean for a system to be secure?

cannot be misused by insiders?

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### **Motivation**

What does it mean for a system to be secure?

only authorized persons have access? only authorized user accounts have access?

What does it mean for a system to be secure?

no buffer overflow bugs?
no buffer overflow vulnerabilities?

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### **Motivation**

What do all of these examples have in common?

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## **POLICY**

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### **Motivation**

What does it mean for a system to be secure?

physically secure?

policy defines...
the physical requirements of the system

What does it mean for a system to be secure?

cannot be misused by insiders?

policy defines...
how the system is *intended* to be used

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### **Motivation**

What does it mean for a system to be secure?

only authorized persons have access? only authorized user accounts have access?

policy defines... who is authorized for what type of access

What does it mean for a system to be secure?

no buffer overflow bugs?
no buffer overflow vulnerabilities?

policy defines...
the difference between bug & vulnerability

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### **Motivation**

What does it mean for a system to be secure?

#### no vulnerabilities

where a *vulnerability* is a set of conditions that may lead to a potential policy violation

How do we define policy?

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## Background

### Background

How do we define policy?

#### **Unifying Policy Hierarchy**

(Adam Carlson, Master's Thesis)

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### **Unifying Policy Hierarchy**

#### **Oracle Policy**

- Represents the intent and will of policy makers
- May not be explicitly specified

#### Example:

Xander is authorized to read file readme.txt

#### **Feasible Policy**

- Represents the intent and will of policy makers
- Takes into account the mechanics and available access controls of the system

#### Example:

User account xander is authorized to read file readme.txt

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## **Unifying Policy Hierarchy**

#### **Configured Policy**

- Represents the policy configured on the machine

#### Example:

All user accounts are authorized to read file readme.txt

#### **Actual Policy**

Represents the policy currently in effect on the machine

#### Example:

No user can read file readme.txt (potentially result of denial of service attack)

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## **Unifying Policy Hierarchy**

#### **Oracle Policy**

Captures policy maker's intent

#### **Feasible Policy**

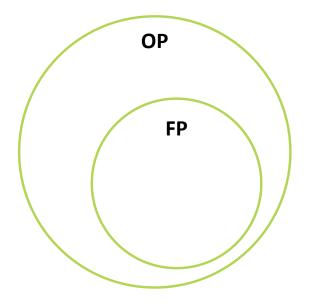
Considers limitations of system

#### **Configured Policy**

Policy as configured on system

#### **Actual Policy**

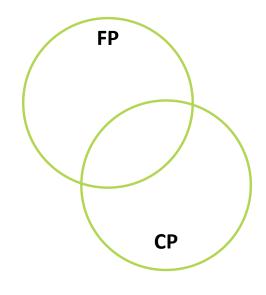
Policy currently in effect on system



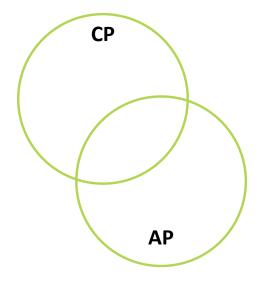
OP ≠ FP Inherent Vulnerability

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## **Unifying Policy Hierarchy**



FP ≠ CP Configuration Vulnerability



CP ≠ AP Runtime Vulnerability

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## **Proposal**

### **Proposal**

1

**Expand application of the hierarchy** 

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## **Proposal**

1 Expand application of the hierarchy

**Insider Threat** 

**Social Engineering** 

**Network Viewpoint** 

And more...

### **Proposal**



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### **Insider Threat**

"exists whenever a lower policy level has *more* authorized privileges than a higher policy level"

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OP: Yasmin may use the system to read medical records to treat patients.

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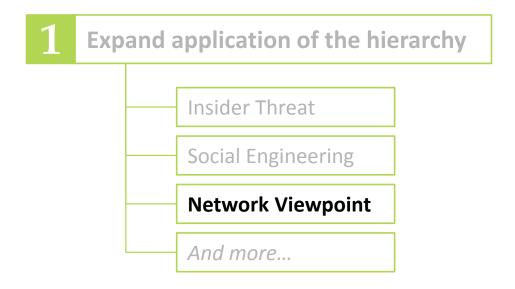
### **Insider Threat**

"exists whenever a lower policy level has more authorized privileges than a higher policy level"

OP: Yasmin may use the system to read medical records to treat patients.

FP: User account yasmin may use the system to read medical records.

### **Proposal**



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## **Network Viewpoint**

In original approach, each system has its own associated policy hierarchy.

### **Network Viewpoint**

In original approach, each system has its own associated policy hierarchy.

How do we expand this to a more network-based approach?

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## **Proposal**

- 1 Expand application of the hierarchy
- 2 Use model to perform threat analysis

## **Threat Analysis**

"Gap Analysis"

Examine the "gap" between levels of the policy hierarchy, i.e. everywhere two consecutive levels do not match.

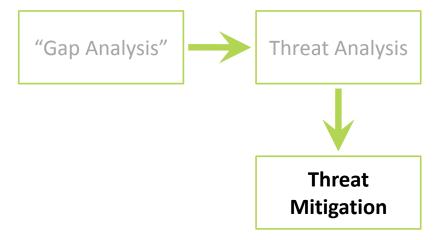
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### **Threat Analysis**



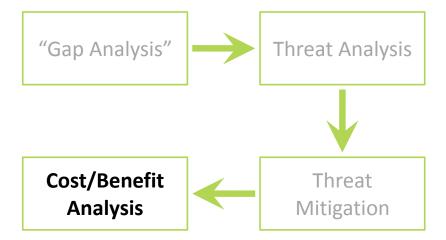
Next, determine the potential threat caused by these gaps.

### **Threat Analysis**

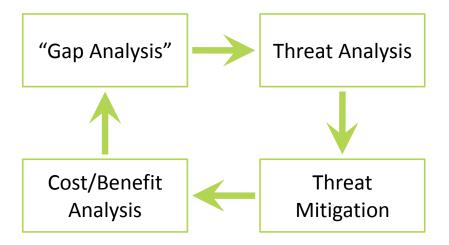


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### **Threat Analysis**



## **Threat Analysis**



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## **Proposal**

- 1 Expand application of the hierarchy
- 2 Use model to perform threat analysis
- 3 Present findings in a wiki format

